# Mathematics for Computer Science: Homework 8

Xingyu Su 2015010697

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### Special Problem 1

In class we discussed the *Mr. P and Mr. S problem*, in which a conversation with 4 rounds of communications takes place. Let  $2 \le m < 98n \le 99$ . Of the  $\binom{98}{2} = 4753$  pairs of (m, n) in the range, let  $a_j$  be the number of pairs leading to this conversation (exactly as stated in class) stopping (and unable to continue) after exactly *j* rounds for j = 0, 1, 2, 3, 4, respectively. You should get  $a_4 = 1$ .

*Mr. P* and *Mr. S* problem: Two numbers *m* and *n* are chosen such that  $2 \le m \le n \le 99$ . *Mr. S* is told their sum and Mr. *P* is told their product. The following dialogue ensues:

Mr. P: I don't know the numbers.

Mr. S: I knew you didn't know. I don't know either.

Mr. P: Now I know the numbers.

Mr. S: Now I know them too.

**Question**: Write a computer program and determine the values of  $a_0, a_1, a_2, a_3, a_4$ . You should give a concise explanation of the principles of your design.

### Answer:

It's not difficult to analyze this problem. Step by step we can get all  $a_j$ .

(1) To get  $a_0$ , we need to know exactly when will *Mr. P don't know the numbers.* That means, the product given to him is not unique of products in all pertubations. Define the whole produts array as *muls*. Count them one by one, we will know the times each *mul* shows up. Then select all situations that *mul* more than one time, we get  $a_0$  situations left.

(2) To get  $a_1$ , we need to know when Mr. S knew Mr. P didn't know. Define the whole sums array as sums. All situations that mul shows exactly one time will make Mr. P know the numbers. So Mr. S knows that the sum given to him will not lead to situation like that. Which means, his sum does not belong to the sums that corresponding muls are unique. Here we get  $a_1$ .

(3) To get  $a_2$ , we need to know when *Mr. S still don't know the numbers*. First we should reduce the scope of results in to  $a_1$ . Then similar to (1), counting all sums and the ones not unique are which the numbers could belong to.  $a_2$  is the number of these situations.

(4)  $a_3$  is the number of situations that *Mr*. *P* know the results under all situations represented by  $a_2$ . So the product *Mr*. *P* got is unique in  $a_2$  situations. Counting this, we get  $a_3$  now.

(5) Mr. S now knows the results lead to similar counting process in (4). Counting all sums in  $a_3$  and find all unique ones. We finally get  $a_4$ .

The program is clear now. I choose *python* as programming language since it's easy to do counting and selection. To get counting part seperated, I used *collections.Counter* function and mapping to write a *IndexFilter* that can get all index that corresponding values in *fvals* satisifying *condition* of counting *cvals*.

The code I wrote is as below:

```
1 #/usr/bin/env python
2 import numpy as np
  from collections import Counter
3
4
   def IndexFilter(cvals, fvals, condition=lambda fvi: True):
5
6
       counter = Counter(cvals)
7
       return [i for i, fvi in enumerate(fvals) if condition(counter[fvi])]
8
9
   if __name__ == '__main__':
10
       # initialize
       numbers = np.array([(m,n) for m in range(2,100) for n in range(m+1,100)])
11
12
       sums = np.array([m+n for (m,n) in numbers])
       muls = np.array([m*n for (m,n) in numbers])
13
14
       # Subsets: when having no infomation, Mr.P don't know
15
       idx_a0 = IndexFilter(muls, muls, condition=lambda fmi: fmi>1)
16
       idx_a0_op = IndexFilter(muls, muls, condition=lambda fmi: fmi==1)
17
18
       # Subsets: Mr.S know Mr.P don't know
19
       opps = sums[idx_a0_op]
20
       idx_a1 = IndexFilter(opps, sums, condition=lambda fsi: fsi==0)
21
22
       # Subsets: Mr.S still don't know.
23
24
       sums,muls = sums[idx_a1],muls[idx_a1]
25
       idx_a2 = IndexFilter(sums, sums, condition=lambda fsi: fsi>1)
26
       # Subsets: Mr.P now know the answer.
27
28
       sums,muls = sums[idx_a2],muls[idx_a2]
29
       idx_a3 = IndexFilter(muls, muls, condition=lambda fmi: fmi==1)
30
       # Subsets: Mr.S now know the answer.
31
32
       sums,muls = sums[idx_a3],muls[idx_a3]
       idx_a4 = IndexFilter(sums, sums, condition=lambda fsi: fsi==1)
33
34
35
       # show results
36
       m,n = numbers[idx_a1][idx_a2][idx_a3][idx_a4][0]
37
       print("a0 = \%d"\%len(idx_a0))
38
       print("a1 = %d"%len(idx_a1))
       print("a2 = \%d"\%len(idx_a2))
39
40
       print("a3 = \%d"\%len(idx_a3))
       print("a4 = \%d"\%len(idx_a4))
41
42
       print("m = %d, n = %d"%(m,n))
```

The results of  $a_i$  are:

 $a_0 = 3021, a_1 = 145, a_2 = 145, a_3 = 86, a_4 = 1.$ 

and the numbers are:

$$m = 4, n = 13, m + n = 17, m \times n = 52.$$

## Special Problem 2

Let C be the unit circle on the complex plane, traversed counter-clockwise. Let n be any integer (positive, negative and zero), and  $A_n = \oint_C z^n$ . C Determine the value of  $A_n$ . You should give a justification of your answer.

### Answer:

To solve the integral along with unit circle on the complex plane, define the unit circle as  $\gamma(\theta) = e^{i\theta}$ . So:

$$A_n = \oint_{\mathcal{C}} z^n = \int_0^{2\pi} e^{in\theta} de^{i\theta}$$
  
= 
$$\int_0^{2\pi} i e^{(n+1)\theta} d\theta$$
  
= 
$$i \int_0^{2\pi} \cos[(n+1)\theta] d\theta - \int_0^{2\pi} \sin[(n+1)\theta] d\theta$$

Easy to know that when  $n+1 \neq 0$ ,  $\int_0^{2\pi} \cos[(n+1)\theta] d\theta = \int_0^{2\pi} \sin[(n+1)\theta] d\theta = 0$ . So  $A_n = 0$ , for  $n \neq -1$ . When n = -1, we have:

$$A_n = \int_0^{2\pi} i d\theta = 2\pi i$$

So, the results are:

$$A_n = \begin{cases} 2\pi i, & n = -1; \\ 0, & else. \end{cases}$$

For justification, consider the Cauchy integral theorem that:

When f is a holomorphic function, and let C be a rectifiable path whose start point is equal to its end point. Then

$$\oint_{\mathcal{C}} f(z) dz = 0.$$

For all  $n \ge 0$ ,  $f(z) = z^n$  is obvious holomorphic, so  $A_n = 0$  for  $n = 0, 1, 2, \cdots$ . The method used before for solving n = -1 is totally calculus and is unnecessary to be discussed here. For the cases that n < -1. We know the Residue theorem that:

$$\oint_{\mathcal{C}} f(z)dz = 2\pi i \sum_{k=1}^{n} \operatorname{Res}(f, z_k)$$

where C is a positively oriented simple closed curve and  $z_k$  are poles. And the residue:

$$Res(f, z_k) = \lim_{z \to z_k} \frac{1}{(k-1)!} \frac{d^{k-1}}{dz^{k-1}} (z-c)^k f(z)$$

So for  $f(z) = z^{-n}$ , n > 1, we have an n order pole z = 0.

$$A_n = \oint_{\mathcal{C}} f(z) dz = 2\pi i \lim_{z \to 0} \frac{1}{(n-1)!} \frac{d^{n-1}}{dz^{n-1}} z^n z^{-n} = 0.$$

In summary, the results still are:

$$A_n = \begin{cases} 2\pi i, & n = -1; \\ 0, & else. \end{cases}$$

#### Acknowledgement:

Thanks to Cauchy integral theorem and Residue theorem for SP2.